

Internal **relational** parametricity, without an interval

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Parametricity

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$$f^P A (\lambda x.x = a) a \text{ refl} : f A a = a$$

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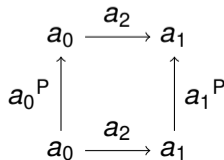
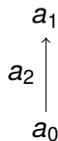
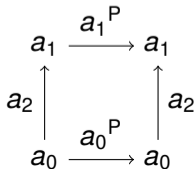
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Dealing with higher dimensional cubes

paper	substructural interval	model	rel/span
Bernardy–Moulin (2012)	no		rel
Bernardy–Moulin (2013)	yes		rel
Bernardy–Coquand–Moulin (2015) Reboullet (2024)	yes	Reedy fibrant presheaf	rel
Nuyts–Devriese (2018)	yes	ordinary presheaf	rel
Cavallo–Harper (2021)	yes	ordinary presheaf	rel
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► Relation vs. span

$A^P : A \rightarrow A \rightarrow \text{Type}$

$\forall A : \text{Type}$ together with $A \xleftarrow{0_A} \forall A \xrightarrow{1_A} A$

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► In this talk we define a relation-based variant of the ACKS theory.

An excerpt of the rules

$$\frac{B : \text{Type}}{B^{\text{P}} : B \rightarrow B \rightarrow \text{Type}}$$

$$\frac{A : B \rightarrow \text{Type} \quad b_2 : B^{\text{P}} b_0 b_1}{\text{Br}_A b_2 : A b_0 \rightarrow A b_1 \rightarrow \text{Type}}$$

An excerpt of the rules

$$\frac{B : \text{Type}}{\forall B : \text{Type} \\ k_B : \forall B \rightarrow B}$$

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- ▶ like Bernardy–Jansson–Paterson (2010)
- ▶ like cubical type theory: “ $\forall B = \mathbb{I} \rightarrow B$ ”

An excerpt of the rules

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$$\overline{\forall (\Sigma B A) = \Sigma (b_2 : \forall B). \Sigma (a_0 : A (0 b_2)). \Sigma (a_1 : A (1 b_2)). \text{Br}_A b_2 a_0 a_1}$$

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$$\frac{A : \text{Type}}{\text{R}_A : A \rightarrow \forall A}$$

$$\frac{A : \text{Type}}{\text{S}_A : \forall (\forall A) \rightarrow \forall (\forall A)}$$

Some equations on k , R , S

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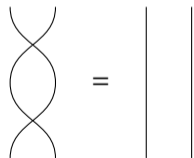
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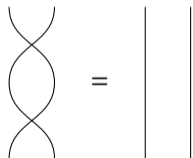


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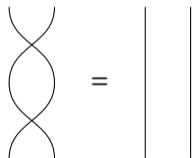
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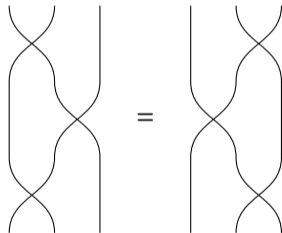
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Computing Bridge

$$\overline{\text{Br}_{A \text{ of } C_2}} = \overline{\text{Br}_A(\text{ap } f \text{ } C_2)}$$

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$$\overline{\text{Br}_{A \circ f} c_2 = \text{Br}_A (\text{ap } f \ c_2)}$$

$$\text{Br}_{\lambda c. (a:A c) \rightarrow B(c,a)} c_2 \ f_0 \ f_1 \cong (a_0 : A (0 \ c_2)) (a_1 : A (1 \ c_2)) (a_2 : \text{Br}_A \ c_2 \ a_0 \ a_1) \rightarrow \text{Br}_B (c_2, a_2) (f_0 \ a_0) (f_1 \ a_1)$$

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$$\text{Br}_{\lambda X.X} : \text{Br}_{\lambda \dots \text{Type}} * A_0 A_1 \leftrightarrow A_0 \rightarrow A_1 \rightarrow \text{Type} : \text{Gel}$$

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$$\overline{\text{gel} : R a_0 a_1 \cong \text{Br}_{\lambda X.X} (\text{Gel } R) a_0 a_1 : \text{ungel}}$$

Polymorphic identity example

Assume

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using the unary version of our theory, we have

$$\text{ungel}(\text{apd } f (A, \text{Gel } P, a, \text{gel } p)) : P (f (A, a)).$$

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using the unary version of our theory, we have

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when setting

$$P := \lambda x. x = a \quad \text{and} \quad p := \text{refl},$$

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- ▶ We use telescopes for \forall , see the abstract for details.

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- ▶ Future work:
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 - ▶ H.O.T.T.: adding transport